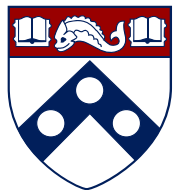


Specifying, Testing and Verifying a Networked Server

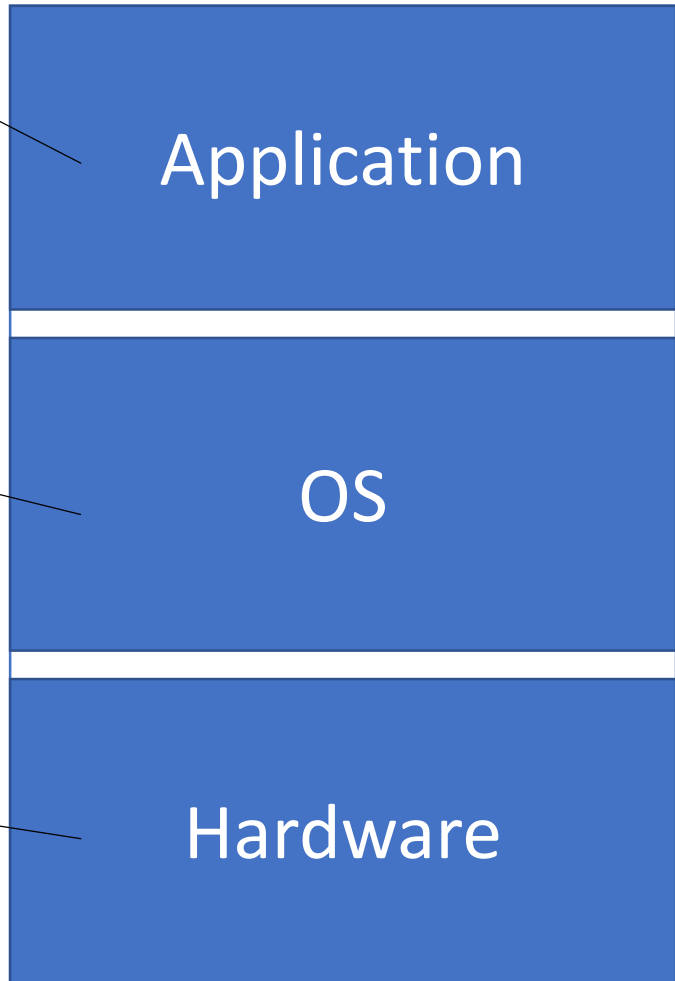
From C to Interaction Trees

Nicolas Koh, Yao Li, Yishuai Li, **Li-yao Xia**
Lennart Beringer, Wolf Honoré, William Mansky
Benjamin C. Pierce, Steve Zdancewic



January 14, 2019 (CPP)

Verification from RFCs to transistors



One theorem to verify...

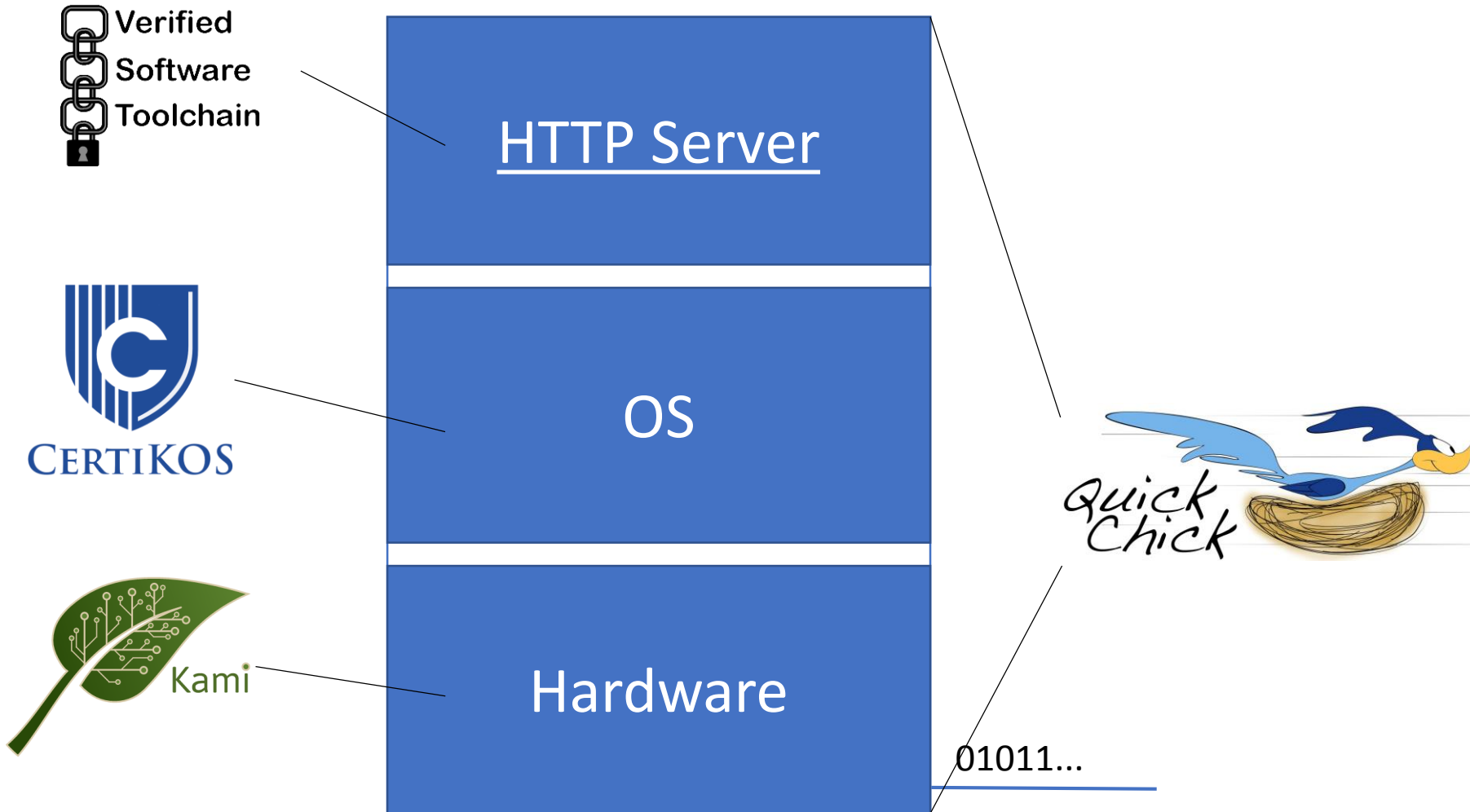


... and test!

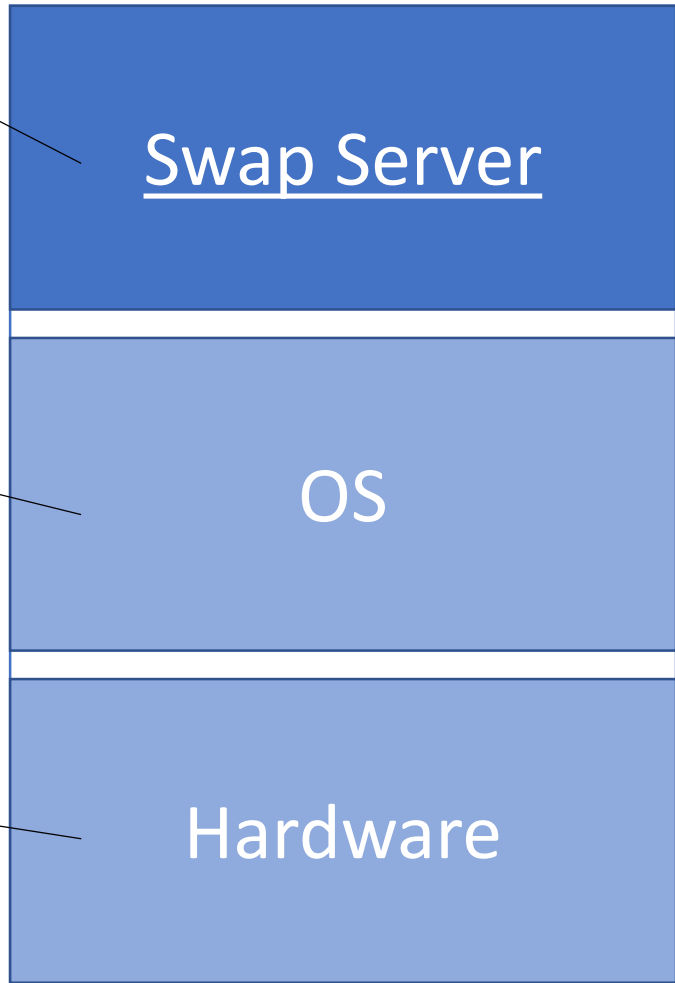
More projects at deepspec.org...

01011...

Towards a verified web server



Towards a verified web server



Today: a simplified server



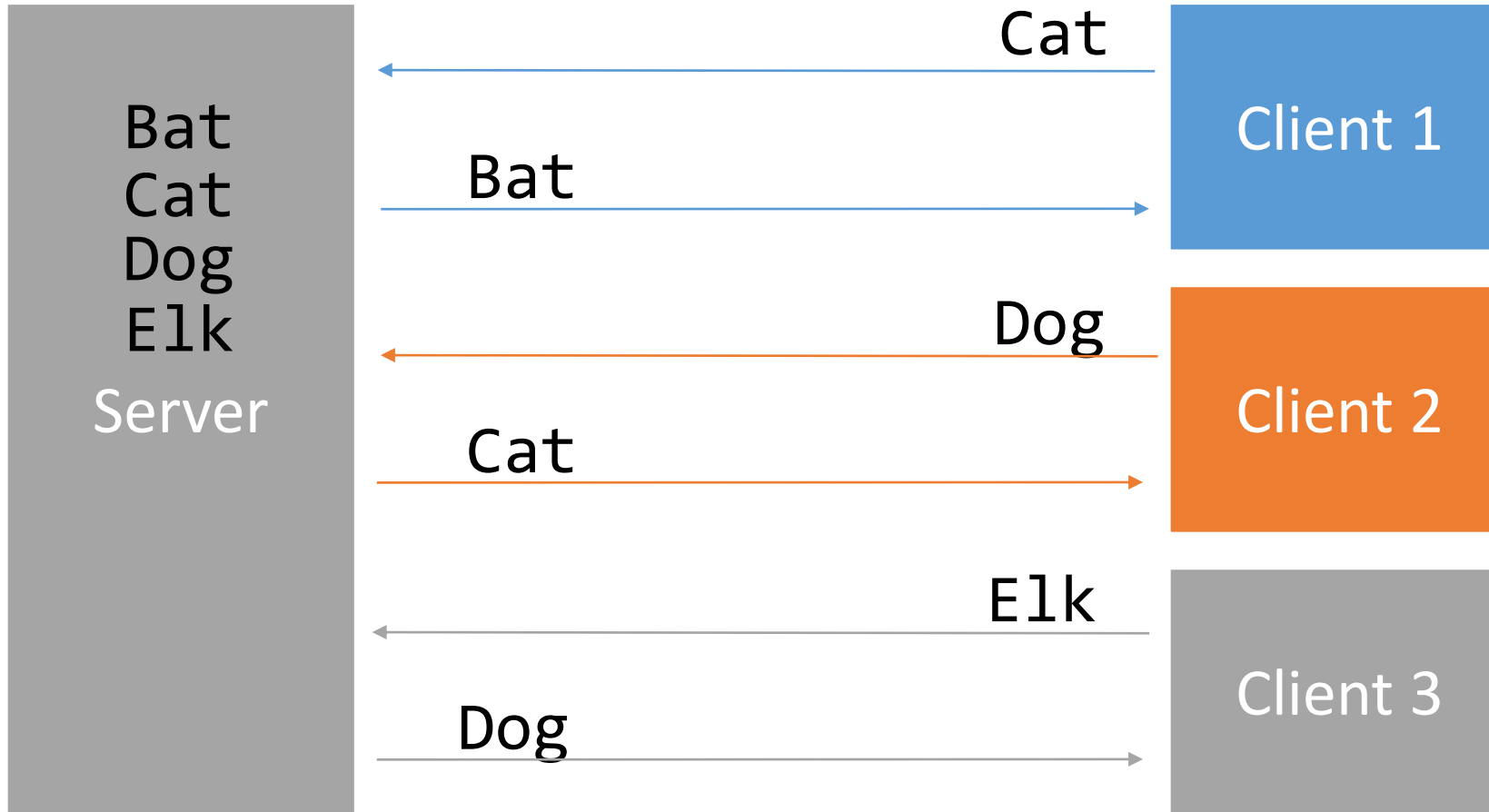
01011...



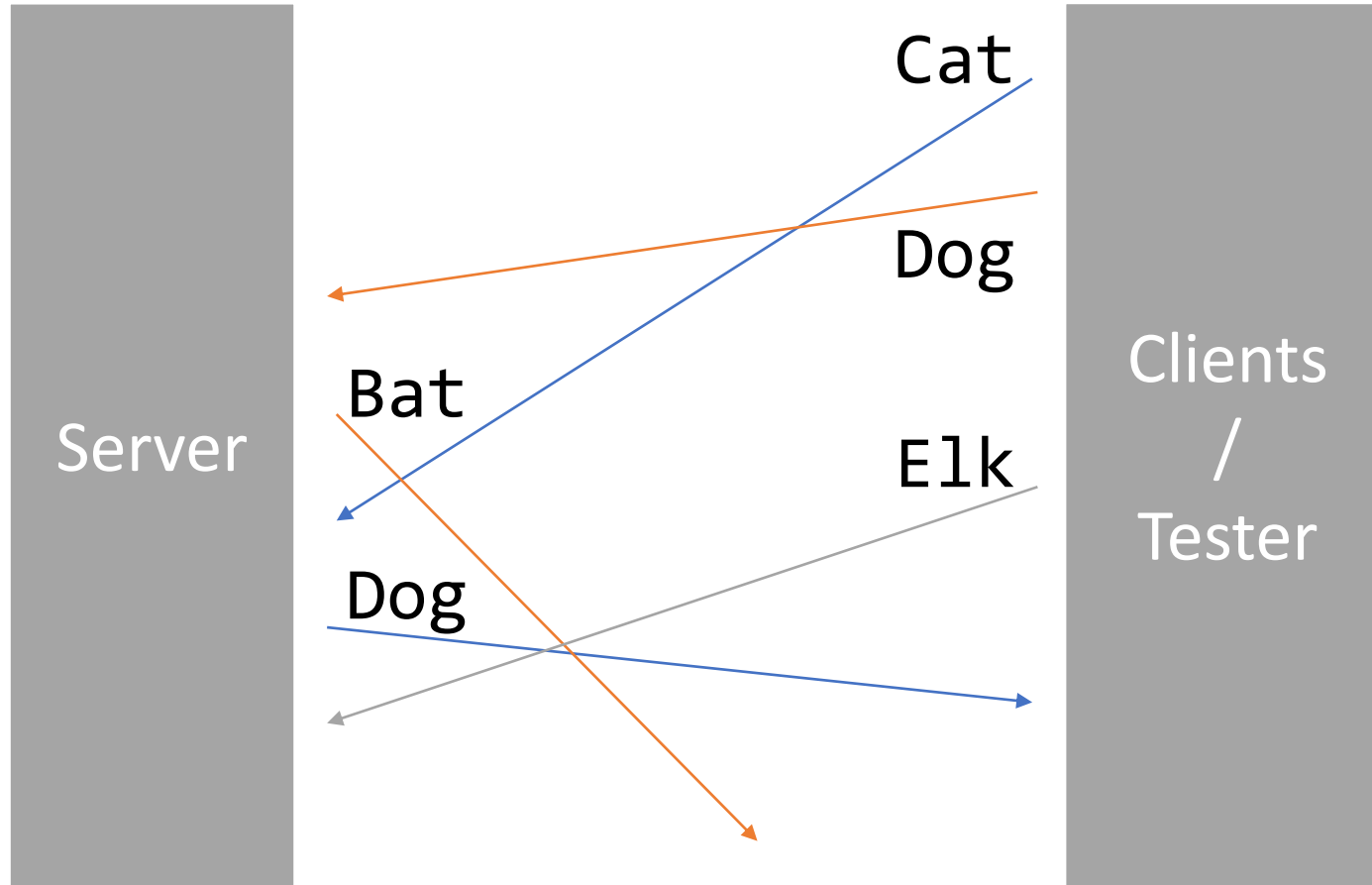
Main contributions

- Verifying a networked C program using VST, which can run in CertiKOS
- Specification describes what a client can observe *over the network*
- Testable specification, using QuickChick

Swap server specification

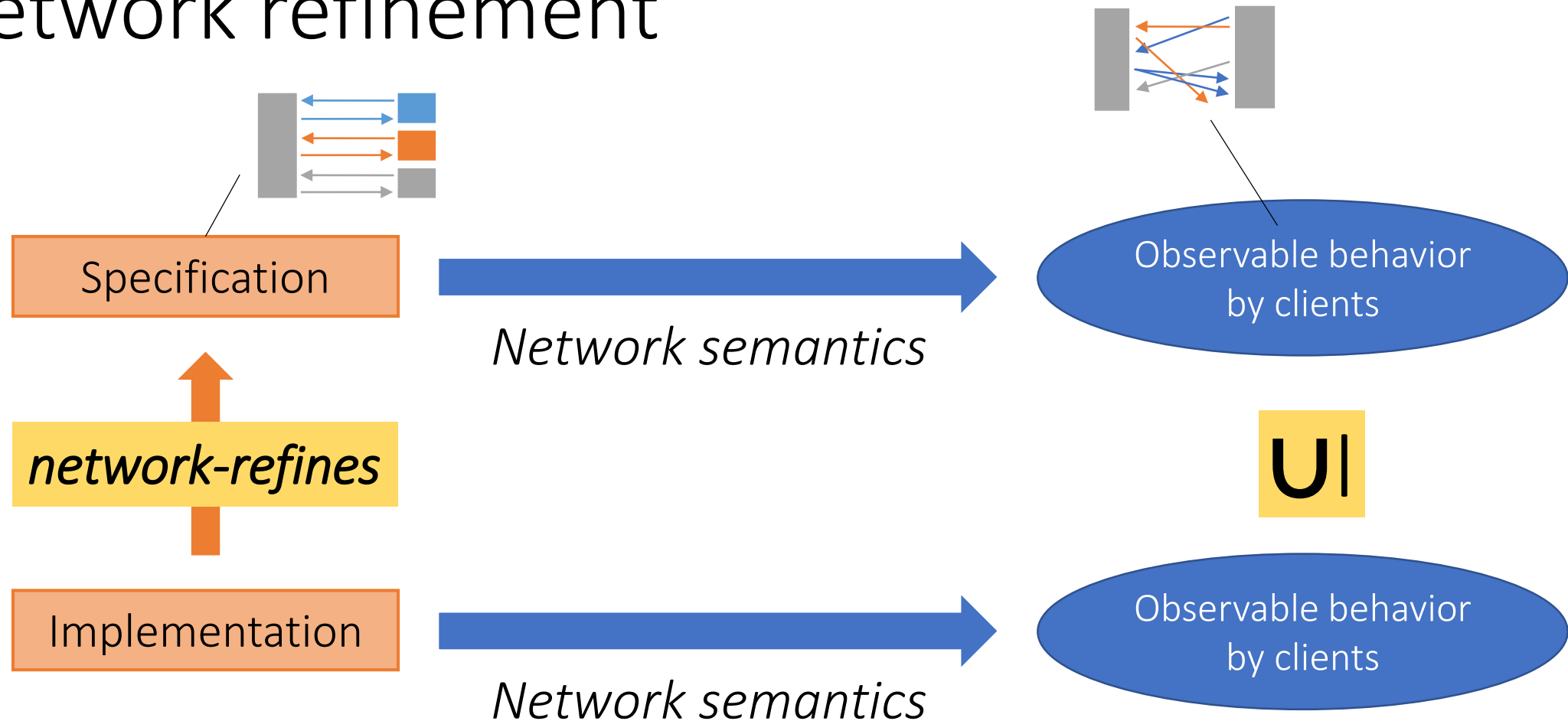


Swap server: in the real world



- Messages on different connections can be reordered
- Messages can be delayed indefinitely

Network refinement



Adaptation of Observational refinement/Linearizability

*: concepts defined in the paper

Overview: proof architecture



Written in Coq

Written in C



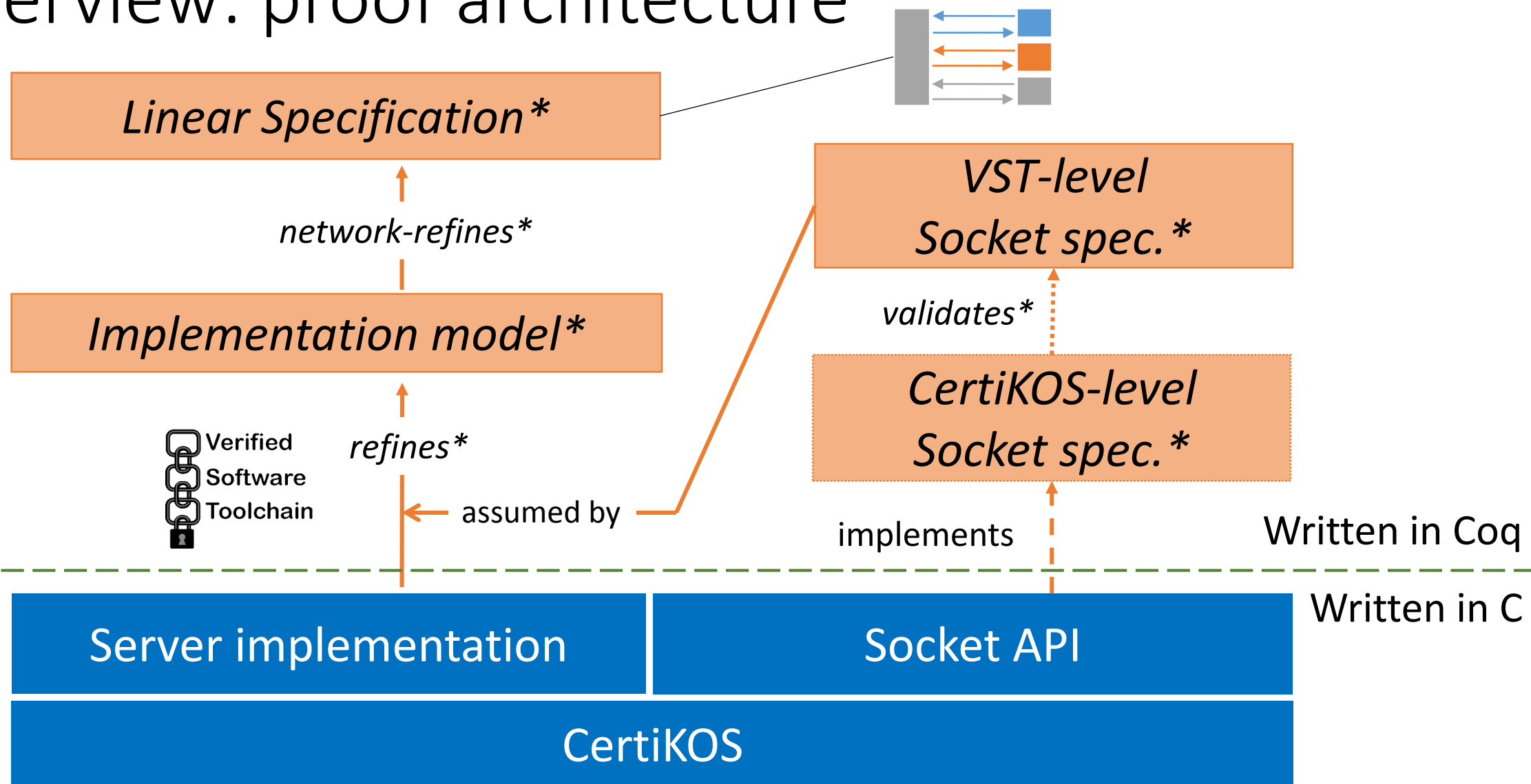
Server implementation

Socket API

CertiKOS

*: concepts defined in the paper

Overview: proof architecture



A unifying specification language

Different spec. styles

Different abstraction levels



testing

*Linear Specification** 

*network-refines**

*Implementation model** 

*VST-level Socket spec.** 

*validates**

*CertiKOS-level Socket spec.** 

*network-refines**



*refines**

assumed by

implements

Written in Coq

Written in C

Server implementation

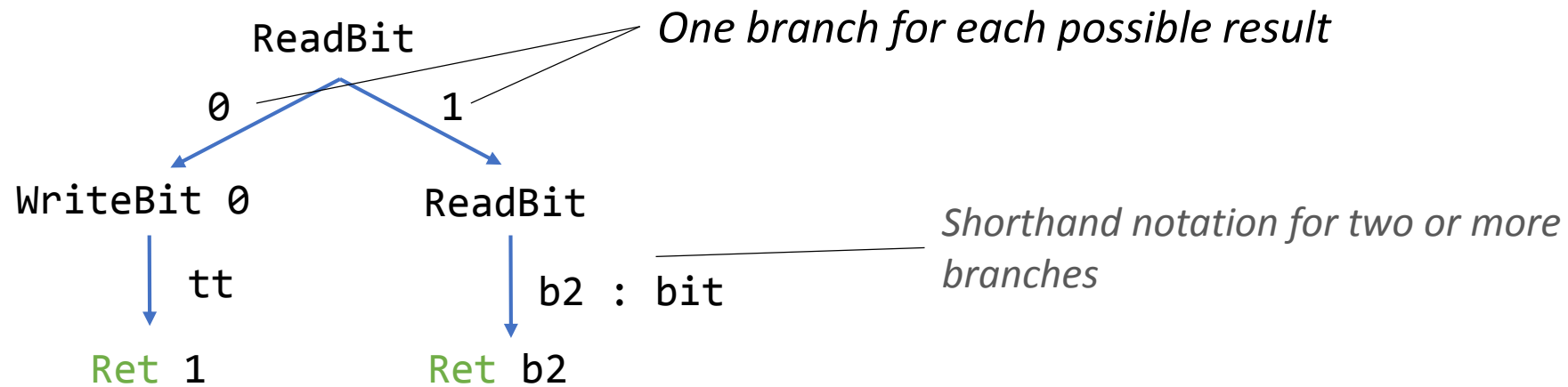
Socket API

CertiKOS



Interaction trees: example

(aka. Free monads)



Type of effects `ioE`:

```
Inductive ioE : Type -> Type :=  
| ReadBit   : ioE bit  
| WriteBit  : bit -> ioE unit  
.
```

Interaction trees: definition

(aka. Free monads)

Type of effects (e.g., `ioE`)

Type of results

CoInductive `itree` (`E` : Type -> Type) (`R` : Type) :
Type :=

| `Vis` : $\forall Y, E\ Y \rightarrow (Y \rightarrow itree\ E\ R) \rightarrow itree\ E\ R$

| `Ret` : $R \rightarrow itree\ E\ R$

| `Tau` : $itree\ E\ R \rightarrow itree\ E\ R$

.

Effect

Continuation

A unifying specification language

Different spec. styles

Different abstraction levels



testing

*Linear Specification** 

*network-refines**

*Implementation model** 

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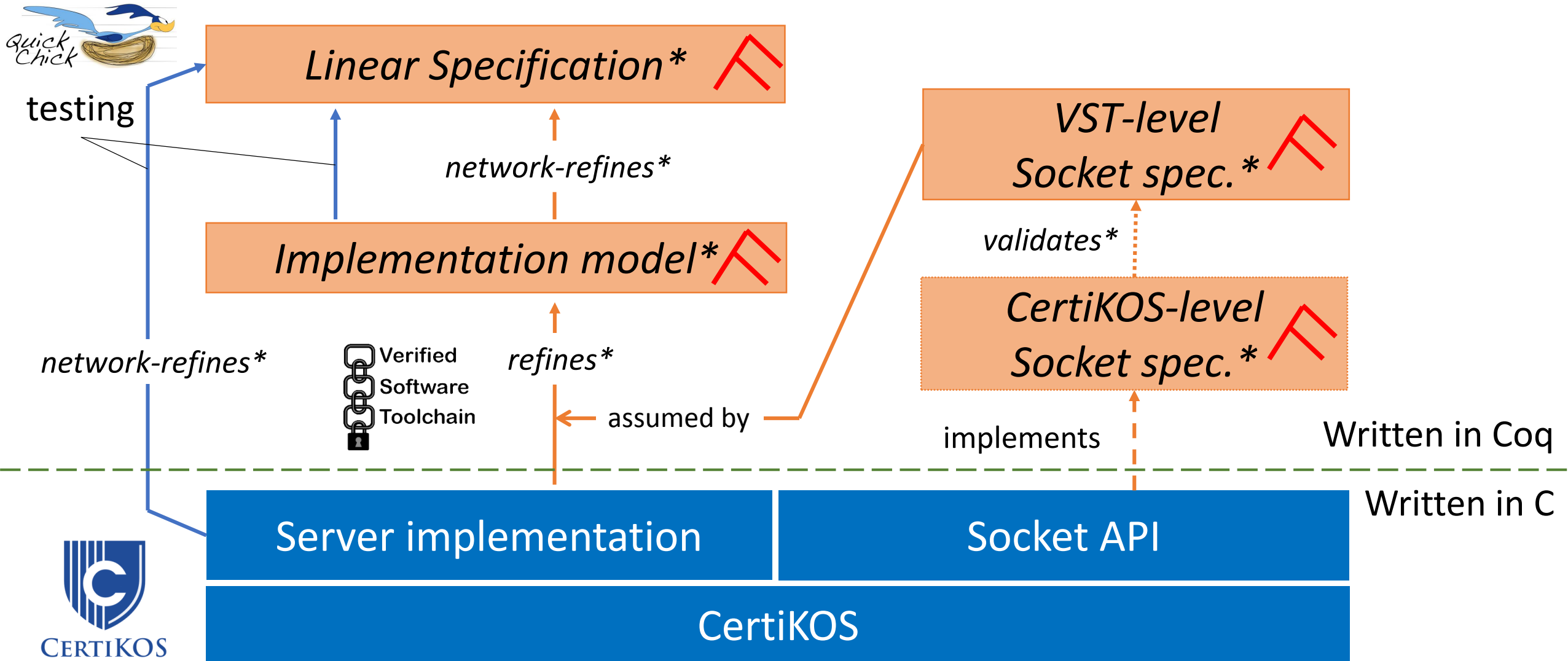
^ The Swap server “linear specification”

CoFixpoint

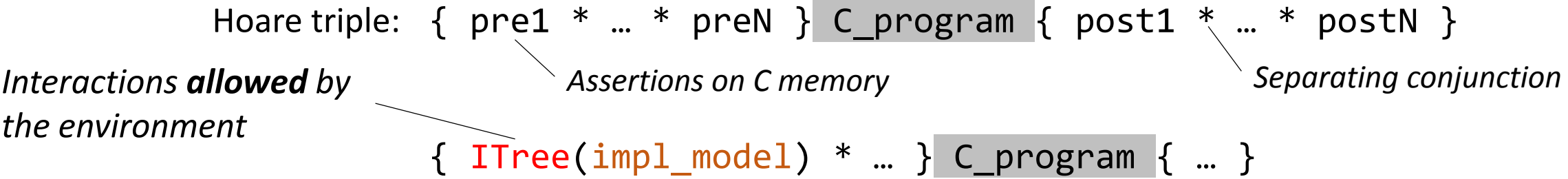
```
loop (open_conns : list conns)
  (last_msg : bytes) : itree serverE unit
:=
  c <- choose open_conns ;;
  new_msg <- recv_msg c ;;
  send_msg c last_msg ;;
  loop open_conns new_msg.
```

Simplified version (see paper)

Overview: proof architecture



Refinement: from C to ITrees



Example of a networked C program with its implementation model:

```

{ ITree(msg <- Recv c ;;
  Send c msg ;;
  t) * ... }
recv(c, buf, len);
send(c, buf, len);
{ ITree(t) * ... }

```

Implementation model (itree)

C implementation

Simplified triples (see paper)

The Swap server correctness theorem

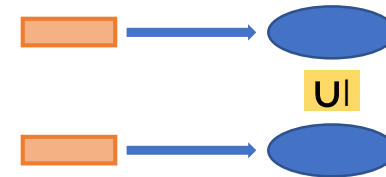
```
{ ITree(impl_model) } C_prog { ... }
```

Theorem correct_server :

exists impl_model,

refines C_prog impl_model /\

network_refines impl_model linear_spec.



Summary and next steps

- Verifying a networked C program using VST, which can run in CertiKOS
- The specification describes a client can observe *over the network*
- The specification is testable, using QuickChick *and Interaction trees*

Scale up:
Swap server ->
HTTP Server

Add more interfaces:
filesystem, encryption...

Improve proof
and testing
techniques

New library

<https://github.com/DeepSpec/InteractionTrees>

